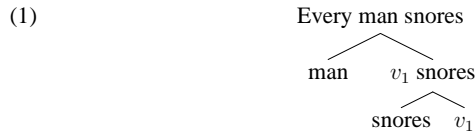
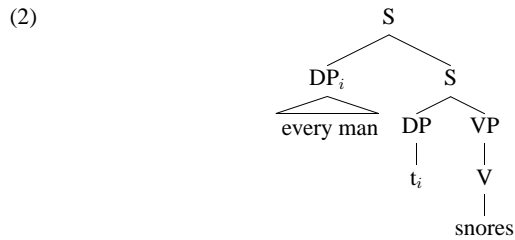


Traces

The Dowty Wall and Peters analysis of quantifiers is very weird:



The Heim and Kratzer one is weird too, but it seems less weird in comparison. As we will see later on, it involves a syntactic transformation of Quantifier Raising (QR), which raises the quantifier *every man* into a higher position in the tree, leaving a trace behind it. **Traces are interpreted as variables.**



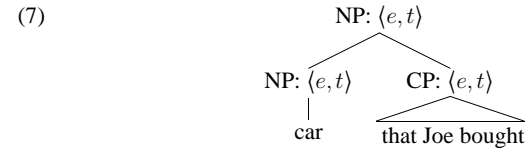
This also works in the analysis of (restrictive) relative clauses, like:

- (3) The car **that Joe bought** is very fancy.
- (4) The woman **who admires Joe** is very lovely.

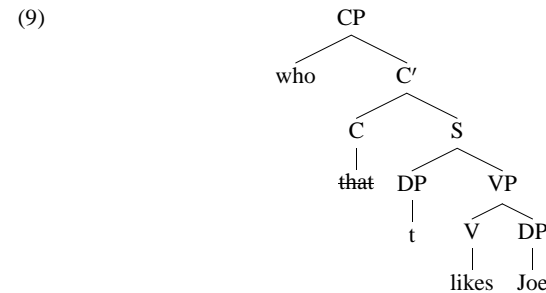
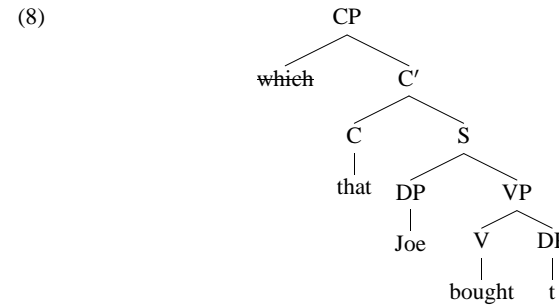
Semantically, relative clauses are just like adjectives:

- (5) The **red** car is very fancy.
- (6) The **Swedish** woman is very lovely.

They are type $\langle e, t \rangle$ and combine via Predicate Modification.



CP stands for “Complementizer Phrase” and Heim and Kratzer assume the following syntax for relative clause CPs:



The text that is struck out like \emptyset is *deleted*. Heim and Kratzer assume that either the relative pronoun *which* or *who* or the complementizer *that* is deleted.

Interpretation of variables Heim and Kratzer use simplified variable assignments. The assignment is just an individual. The interpretation of a trace with respect to this assignment is the individual.

(10) $\llbracket t \rrbracket^{\text{Mary}} = \text{Mary}$

(11) **Traces Rule**

If α is a trace and a is an assignment, $\llbracket \alpha \rrbracket^a = a$

So now we interpret everything with respect to an assignment.

(12) $\llbracket [_{\text{VP}} [_{\text{V}} \text{abandoned}] [_{\text{DP}} t]] \rrbracket^{\text{Mary}} = \lambda x . x \text{ abandoned Mary}$

(13) $\llbracket [_{\text{VP}} [_{\text{V}} \text{abandoned}] [_{\text{DP}} t]] \rrbracket^{\text{Fred}} = \lambda x . x \text{ abandoned Fred}$

But there are *assignment-independent* denotations too.

(14) For any tree α , α is in the domain of $\llbracket \cdot \rrbracket$ iff for all assignments a and b , $\llbracket \alpha \rrbracket^a = \llbracket \alpha \rrbracket^b$.

If α is in the domain of $\llbracket \cdot \rrbracket$, then for all assignments a , $\llbracket \alpha \rrbracket = \llbracket \alpha \rrbracket^a$.

So we can still have assignment-independent lexical entries like:

(15) $\llbracket \text{laugh} \rrbracket = \lambda x \in D_e . x \text{ laughs}$

and then by (14), we have:

(16) $\llbracket \text{laugh} \rrbracket^{\text{Mary}} = \lambda x \in D_e . x \text{ laughs}$

(17) $\llbracket \text{laugh} \rrbracket^{\text{Fred}} = \lambda x \in D_e . x \text{ laughs}$

We need to redo the composition rules now too:

(18) **Lexical Terminals**

If α is a terminal node occupied by a lexical item, then $\llbracket \alpha \rrbracket$ is specified in the lexicon.

(19) **Non-branching Nodes (NN)**

If α is a non-branching node and β its daughter, then, for any assignment a , $\llbracket \alpha \rrbracket^a = \llbracket \beta \rrbracket^a$.

(20) **Functional Application (FA)**

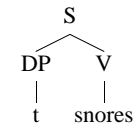
If α is a branching node and $\{\beta, \gamma\}$ the set of its daughters, then, for any assignment a , if $\llbracket \beta \rrbracket^a$ is a function whose domain contains $\llbracket \gamma \rrbracket^a$, then $\llbracket \alpha \rrbracket^a = \llbracket \beta \rrbracket^a(\llbracket \gamma \rrbracket^a)$.

(21) **Predicate Modification (PM)**

If α is a branching node and $\{\beta, \gamma\}$ the set of its daughters, then, for any assignment a , if $\llbracket \beta \rrbracket^a$ and $\llbracket \gamma \rrbracket^a$ are both functions of type $\langle e, t \rangle$, then $\llbracket \alpha \rrbracket^a = \lambda x \in D . \llbracket \beta \rrbracket^a(x) = \llbracket \gamma \rrbracket^a(x) = 1$.

Exercise: Compute $\llbracket (22) \rrbracket^{\text{Mary}}$.

(22)

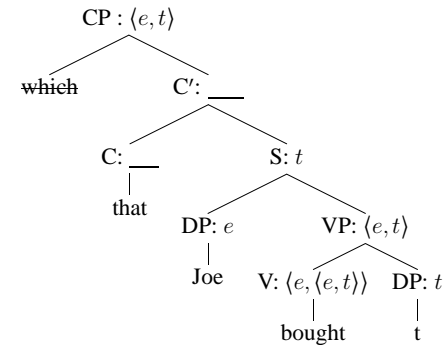


Assume that the people who snore are Mary, John, and Bob.

Exercise: Show that $\llbracket (22) \rrbracket$ is undefined; i.e., (22) is not in the domain of $\llbracket \cdot \rrbracket$.

Predicate abstraction. The S in a relative clause is type t . How do we get the CP to have type $\langle e, t \rangle$?

(23)



Heim and Kratzer:

- The complementizer *that* is vacuous; *that* $S = S$
- The relative pronoun is vacuous too, but it triggers a special rule called Predicate Abstraction

(24) **Predicate Abstraction (PA)**

If α is a branching node whose daughters are a relative pronoun and β , then $\llbracket \alpha \rrbracket = \lambda x \in D_e . \llbracket \beta \rrbracket^x$

So $\llbracket (23) \rrbracket = \lambda x \in D_e . \text{Joe bought } x$.