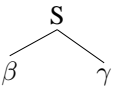




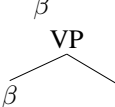


Type-driven interpretation

Syntax-driven interpretation (Heim and Kratzer, ch. 2)

- (S1) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \gamma \rrbracket(\llbracket \beta \rrbracket)$.
- (S2) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.
- (S3) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.
- (S4) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.
- (S5) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.
- (S6) If α has the form  then $\llbracket \alpha \rrbracket = \llbracket \gamma \rrbracket(\llbracket \beta \rrbracket)$.

Type-driven interpretation (Heim and Kratzer pp. 48-9, 65)

- (1) *Terminal Nodes* (TN)
 If α is a terminal node, $\llbracket \alpha \rrbracket$ is specified in the lexicon.
- (2) *Non-Branching Nodes* (NN)
 If α is a non-branching node, and β is its daughter node, then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$.
- (3) *Functional Application* (FA)
 If α is a branching node, $\{\beta, \gamma\}$ is the set of α 's daughters, and $\llbracket \beta \rrbracket$ is a function whose domain contains $\llbracket \gamma \rrbracket$, then $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket(\llbracket \gamma \rrbracket)$.
- (4) *Predicate Modification* (PM)
 If α is a branching node, $\{\beta, \gamma\}$ is the set of α 's daughters, and $\llbracket \beta \rrbracket$ and $\llbracket \gamma \rrbracket$ are both in $D_{(e,t)}$, then
 $\llbracket \alpha \rrbracket = \lambda x \in D_e. \llbracket \beta \rrbracket(x) = 1 \text{ and } \llbracket \gamma \rrbracket(x) = 1$