Lecture 7 Problem of Quantifiers in Object Position & Presupposition

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Introduction to Semantics · EGG 2019

Outline

Problem of Quantifiers in Object Position

QR vs. Direct compositionality

Summary

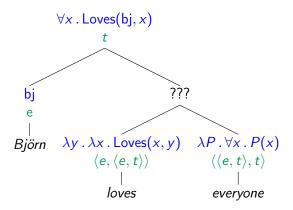
Introducing Presupposition

Presupposition: Formal analysis

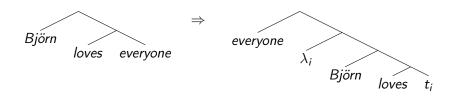
Quantifier in object position

Björn loves everyone $\rightsquigarrow \forall x . \mathsf{Loves}(\mathsf{bj}, x)$

Problem of quantifiers in object position



Quantifier Raising



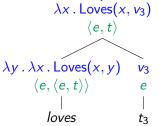
The T-model (Government and Binding Theory)



Quantifier Raising $\forall x$. Loves(bj, x) $\lambda P \cdot \forall x \cdot P(x) \quad \lambda v_3 \cdot \text{Loves(bj, } v_3)$ $\langle\langle e,t\rangle,t\rangle$ $\langle e, t \rangle$ everyone Loves(bj, v_3) λ_3

bį

Björn



 $\langle e, \langle e, t \rangle \rangle$

loves

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Direct compositionality

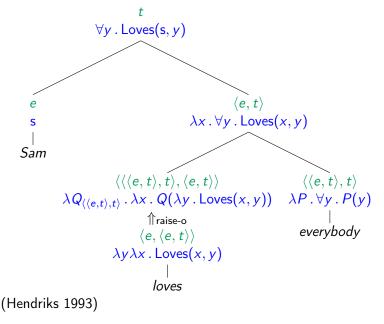
Direct compositionality (in slogan form)

The syntax and semantics work in tandem.

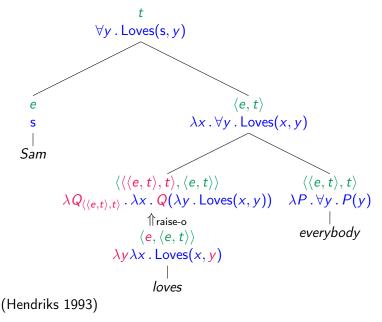
⇒ Each expression computed by the syntax can be interpreted; interpretation is not 'postponed' to a later stage (Jacobson 2012, i.a.).

Does the problem of quantifiers in object position require a violation of direct compositionality?

A directly compositional approach: Type-shifting the verb



A directly compositional approach: Type-shifting the verb



Is QR empirically motivated?

Arguments by Heim & Kratzer in favor of QR against a (different) type-shifting analysis:

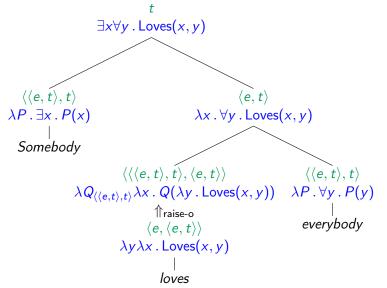
- Scope ambiguities
- Inverse linking
- Antecedent-contained deletion
- Quantifiers that bind pronouns
- Extraction-scope generalization

Scope ambiguities: the problem

Scope ambiguities: the problem

- (1) Somebody loves everybody.
 - ▶ $\forall > \exists$: For every person y: there is a person x such that x loves y.
 - ∃ > ∀
 There is a person x such that for all y: x loves y.

For surface scope: Shift the object, then combine with the subject.



For inverse scope: Shift the subject first, then the object.

```
\forall y \exists x . Loves(x, y)
       \langle\langle e, t \rangle, t \rangle
                                                                                                    \langle\langle\langle e, t \rangle, t \rangle, t \rangle
                                                                    \lambda Q_{\langle\langle e,t\rangle,t\rangle}. \forall y. Q(\lambda x. Loves(x,y))
  \lambda P \cdot \exists x \cdot P(x)
      Somebody
                                 \langle\langle\langle e, t \rangle, t \rangle, \langle\langle\langle e, t \rangle, t \rangle, t \rangle\rangle
                                                                                                                                                         \langle\langle e,t\rangle,t\rangle
\lambda Q'_{\langle\langle e,t\rangle,t\rangle} \lambda Q_{\langle\langle e,t\rangle,t\rangle} \cdot Q'(\lambda y \cdot Q(\lambda x \cdot \mathsf{Loves}(x,y))) \quad \lambda P \cdot \forall y \cdot P(y)
                                                            ↑raise-o
                                                                                                                                                         everybody
                                             \langle e, \langle \langle \langle e, t \rangle, t \rangle, t \rangle \rangle
                      \lambda y \lambda Q_{\langle\langle e,t \rangle,t \rangle}. Q(\lambda x \cdot \mathsf{Loves}(x,y))
                                                            ↑raise-s
                                                         \langle e, \langle e, t \rangle \rangle
                                            \lambda y \lambda x. Loves(x, y)
                                                               loves
```

For inverse scope: Shift the subject first, then the object.

```
\forall y \exists x . Loves(x, y)
       \langle\langle e, t \rangle, t \rangle
                                                                                                   \langle\langle\langle e, t \rangle, t \rangle, t \rangle
                                                                   \lambda Q_{\langle\langle e,t\rangle,t\rangle}. \forall y. Q(\lambda x. Loves(x,y))
  \lambda P \cdot \exists x \cdot P(x)
      Somebody
                                 \langle\langle\langle e, t \rangle, t \rangle, \langle\langle\langle e, t \rangle, t \rangle, t \rangle\rangle
                                                                                                                                                       \langle\langle e,t\rangle,t\rangle
\lambda Q'_{(\langle e,t\rangle,t\rangle} \lambda Q_{(\langle e,t\rangle,t\rangle} \cdot Q'(\lambda y \cdot Q(\lambda x \cdot \mathsf{Loves}(x,y))) \quad \lambda P \cdot \forall y \cdot P(y)
                                                           ↑raise-o
                                                                                                                                                       everybody
                                             \langle e, \langle \langle \langle e, t \rangle, t \rangle, t \rangle \rangle
                      \lambda y \lambda Q_{\langle\langle e,t \rangle,t \rangle}. Q(\lambda x \cdot \mathsf{Loves}(x,y))
                                                           ↑raise-s
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                                                              loves
```

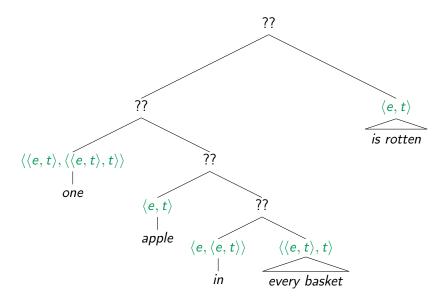
For inverse scope: Shift the subject first, then the object.

```
\forall y \exists x . Loves(x, y)
       \langle\langle e, t \rangle, t \rangle
                                                                                                  \langle\langle\langle e, t \rangle, t \rangle, t \rangle
                                                                   \lambda Q_{\langle\langle e,t\rangle,t\rangle}. \forall y. Q(\lambda x. Loves(x,y))
  \lambda P \cdot \exists x \cdot P(x)
      Somebody
                                 \langle\langle\langle e, t \rangle, t \rangle, \langle\langle\langle e, t \rangle, t \rangle, t \rangle\rangle
                                                                                                                                                     \langle\langle e,t\rangle,t\rangle
\lambda Q'_{\langle (e,t),t \rangle} \lambda Q_{\langle (e,t),t \rangle} \cdot Q'(\lambda y \cdot Q(\lambda x \cdot \mathsf{Loves}(x,y))) \quad \lambda P \cdot \forall y \cdot P(y)
                                                           ↑raise-o
                                                                                                                                                     everybody
                                            \langle e, \langle \langle \langle e, t \rangle, t \rangle, t \rangle \rangle
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                                                              loves
```

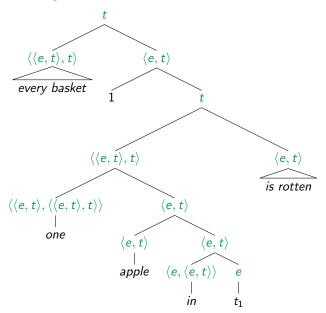
Inverse linking: the problem

- (2) One apple in every basket is rotten.
 - ∀ > ∃:
 For every basket y: there is an apple x in y that is rotten.
 - ▶ $\exists > \forall$ (unavailable) There is an apple that is in every basket and also rotten.

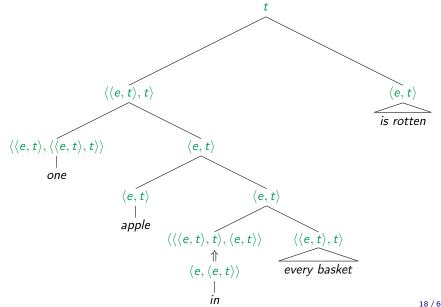
Inverse linking: the problem



Inverse linking: QR solution



Inverse linking: What happens if we object-raise in?



Inverse linking: Directly compositional solution

See Barker 2005: 'Remark on Jacobson 1999: Crossover as a local constraint', *Linguistics and Philosophy*.

Antecedent-contained deletion

VP-ellipsis:

(3) I read War and Peace before you did read War and Peace.

Antecedent-contained deletion with and without QR:

- (4) a. Mary read every novel that John did read every novel that John did read every
 - b. [Every novel that John did read t] Mary read t

Antecedent-contained deletion

VP-ellipsis:

(3) I read War and Peace before you did read War and Peace.

Antecedent-contained deletion with and without QR:

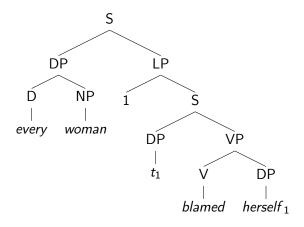
- (4) a. Mary read every novel that John did read every novel that John did read every
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Jacobson (1999) solves this without QR, using function composition.

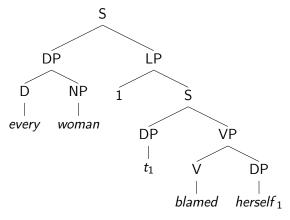
Quantifiers that bind pronouns

- (5) a. Mary blamed herself.
 - b. Mary blamed Mary.
- (6) a. Every woman blamed herself.
 - b. Every woman blamed every woman.
- (7) a. No man noticed the snake next to him.
 - b. No man noticed the snake next to no man.

Quantifiers that bind pronouns: QR solution



Quantifiers that bind pronouns: QR solution



Without QR, herself 1 would end up as a free variable:

$$\forall x . [\mathsf{Woman}(x) \to \mathsf{Blamed}(x, v_1)]$$

Quantifiers that bind pronouns without QR

There are non-QR options, including:

(8) herself
$$\rightsquigarrow \lambda R_{\langle e, \langle e, t \rangle \rangle} \lambda x \cdot R(x, x)$$

Related: Strict and sloppy readings

From *Ghostbusters*:

(9) Dr Ray Stantz: You know, it just occurred to me that we really haven't had a successful test of this equipment.

Dr. Egon Spengler: I blame myself.

Dr. Peter Venkman: So do I.

strict reading: Venkman blames Spengler sloppy reading: Venkman blames Venkman

Related: Strict and sloppy readings

From *Ghostbusters*:

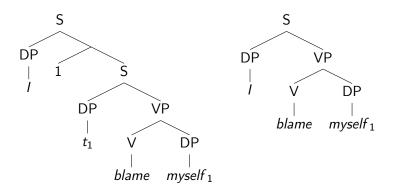
(9) Dr Ray Stantz: You know, it just occurred to me that we really haven't had a successful test of this equipment.

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strict reading: Venkman blames Spengler \leftarrow free pronoun **sloppy reading**: Venkman blames Venkman \leftarrow bound pronoun

Strict and sloppy readings with QR



But Polly Jacobson has addressed this issue as well.

Extraction-scope generalization

In many cases, the constraints on *wh*-extraction mirror the constraints on scope.

- (10) a. John knows a woman from every country.
 - b. #John knows a woman who is from every country.
- (11) a. Which country does John know a woman from?
 - b. *Which country does John know a woman who is from?

But many counterexamples; cf. Simon Charlow's lecture last Wednesday.

So do we really need QR?

No, direct compositionality can be maintained (as far as I can see).

You just need a bit of advanced machinery to handle:

- Antecedent-contained deletion
- Inverse linking
- Strict vs. sloppy identity

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QR vs. Direct compositionality

Summary

Introducing Presupposition

Presupposition: Formal analysis

Summary

Composition rules:

- ► Function Application
- Predicate Modification
- Pronouns and Traces Rule
- Predicate Abstraction

Composition Rules (I)

Function Application

Let γ be a tree whose only two subtrees are α and β where:

- $ightharpoonup \alpha \leadsto \alpha'$ and α' has type $\langle \sigma, \tau \rangle$
- $\triangleright \beta \leadsto \beta'$ and β' has type σ .

Then

$$\gamma \leadsto \alpha'(\beta')$$

Composition Rules (II)

Predicate Modification

lf:

- $ightharpoonup \gamma$ is a tree whose only two subtrees are α and β
- $\triangleright \alpha \leadsto \alpha'$
- $\triangleright \beta \leadsto \beta'$
- $ightharpoonup \alpha'$ and β' are of type $\langle e, t \rangle$

Then:

$$\gamma \rightsquigarrow \lambda u . [\alpha'(u) \land \beta'(u)]$$

where u is a variable of type e that does not occur free in α' or β' .

Composition rules (III)

Pronouns and Traces Rule

If α is an indexed trace or pronoun, $\alpha_i \rightsquigarrow v_{e,i}$

Predicate Abstraction

lf

- $\triangleright \gamma$ is an expression whose only two subtrees are α_i and β
- $\triangleright \beta \leadsto \beta'$
- $\triangleright \beta'$ is an expression of type t

Then $\gamma \rightsquigarrow \lambda v_{i,e} \cdot \beta'$

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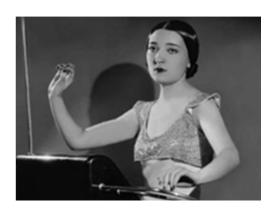
Summary

Introducing Presupposition

Presupposition: Formal analysis

The theremin

- ▶ Electronic instrument
- Controlled without physical contact
- ► Patented in 1928 by Léon Theremin.
- ➤ Clara Rockmore (1911–1998) was a theremin virtuoso



Wolfgang Amadeus Mozart



- Lived 1756-1771
- Composed pieces for many different instruments
- Never encountered a theremin
- As a consequence, never composed a theremin duo

(12) There is at least one theremin duo by Mozart.

(12) There is at least one theremin duo by Mozart. Not true

(12) There is at least one theremin duo by Mozart. Not true

Negation:

(12) There is at least one theremin duo by Mozart. Not true

Negation:

(13) There are no theremin duos by Mozart.

(12)There is at least one theremin duo by Mozart. Not true

Negation:

(13)There are no theremin duos by Mozart. True

(14) Every theremin duo by Mozart is famous.

(14) Every theremin duo by Mozart is famous. Not true

(14) Every theremin duo by Mozart is famous.

Not true

Negation:

(14) Every theremin duo by Mozart is famous.

Not true

Negation:

(15) Not every theremin duo by Mozart is famous.

(14) Every theremin duo by Mozart is famous. Not true

Negation:

(15) Not every theremin duo by Mozart is famous. Not true

Presupposition

If A presupposes B, then A not only implies B but also implies that the truth of B is somehow taken for granted, treated as uncontroversial.

(Chierchia & McConnell-Ginet 2000, 28)

Presupposition accommodation

On Jimmy Kimmel's Lie Witness News:

INTERVIEWER: What do you think of the government's plan to schedule earthquakes for every five years, instead of the current 12 years?

INTERVIEWEE: I think it's very.... conservation-minded.

Presupposition

If A presupposes B, then to assert A, deny A, wonder whether A, or suppose A – to express any of these attitudes toward A is generally to imply B, to suggest that B is true and, moreover, uncontroversially so. That is, considering A from almost any standpoint seems already to assume or presuppose the truth of B; B is part of the background against [which] we (typically) consider A.

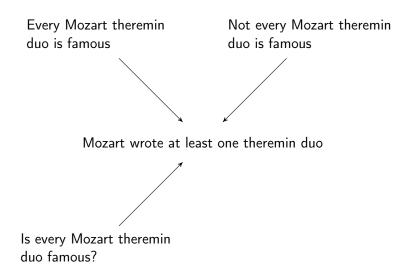
(Chierchia & McConnell-Ginet 2000, 28)

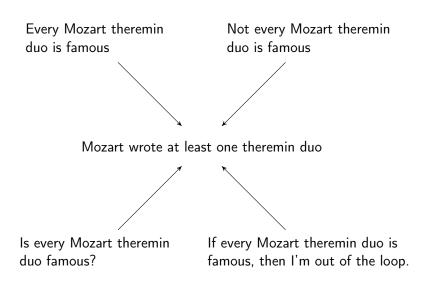
Every Mozart theremin duo is famous

Mozart wrote at least one theremin duo

Every Mozart theremin duo is famous

Mozart wrote at least one theremin duo





Presupposition projection

If A implies B, and 'not A' implies B, then inference from A to B projects over negation.

Presuppositions project (i) over negation, (ii) through question formation, and (iii) from the antecedent of a conditional.

To run the projection test

- 1. Construct the examples:
 - (i) not-A (the negation of A)
 - (ii) A? (a yes/no question)
 - (iii) If A, then C (a conditional)
- Ask the question:Do these sentences imply B?
- 3. Interpret the result: Yes \Rightarrow presupposition.

Triggers

A presupposition trigger is a word or construction that conventionally signals a presupposition.

Example: every.

Presupposition triggers

Ed is glad we won \gg We won Ed knows we won \gg We won Ed's son is bald \gg Ed has a son Only Ed came \gg Ed came The balcony is lovely \gg There is a balcony

To paraphrase Strawson

Your friend says:

(16) The king of France is wise. (using an empty definite description)

Would you agree or disagree? Preferably neither.

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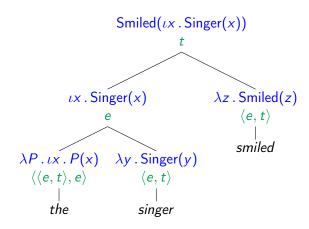
QR vs. Direct compositionality

Summary

Introducing Presupposition

Presupposition: Formal analysis

Definite descriptions (Fregean analysis)



lota

Syntax rule: lota

If ϕ is an expression of type t, and u is a variable of type e, then $\iota u \cdot \phi$ is an expression of type e.

Semantic rule: lota

$$\llbracket \iota \mathbf{u} \, . \, \phi \rrbracket^{M, \mathsf{g}} = \left\{ \begin{array}{l} d \text{ if } \{k : \llbracket \phi \rrbracket^{M, \mathsf{g}[\mathbf{u} \mapsto k]} = 1\} = \{d\} \\ \#_{\mathbf{e}} \text{ otherwise} \end{array} \right.$$

 $\#_e$: 'a completely alien entity' (Kaplan 1989), what definite descriptions denote when they don't denote anything that exists.

Sentences with empty descriptions

If a sentence contains an empty description, what does the sentence denote?

- ► Normally, the third truth value: #
- Exception: the existence predicate
- ▶ Other exceptions: with presupposition plugs and filters

Existence predicate with empty descriptions

- (17) The golden mountain does not exist.
- (18) $\neg \mathsf{Exists}(\iota x . [\mathsf{Golden}(x) \land \mathsf{Mountain}(x)])$

Semantic rule: Existence predicate

 $\llbracket \mathsf{Exists}(lpha)
rbracket^{M,g} = 1$ if $\llbracket lpha
rbracket^{M,g}
eq \#_e$ and 0 otherwise

Other predicates with empty descriptions

- (19) The golden mountain is in Nebraska. $In(\iota x . [Golden(x) \land Mountain(x)], nebraska)$
- (20) The golden mountain is not in Nebraska. $\neg ln(\iota x . [Golden(x) \land Mountain(x)], nebraska)$

Both would normally denote the third truth value, #.

Negation in 3-valued logic



Other presupposition triggers

- (21) a. Both candidates laughed.
 - b. If both candidates laughed, then...
- (22) a. Neither candidate laughed.
 - b. If neither candidate laughed, then...

All imply: There were two candidates.

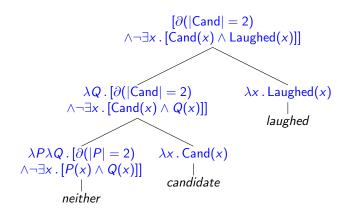
Semantics of ∂ 'partial'



Example

```
Neither candidate laughed \rightsquigarrow [\partial(|\mathsf{Cand}| = 2) \land \neg \exists x . [\mathsf{Cand}(x) \land \mathsf{Laughed}(x)]]
```

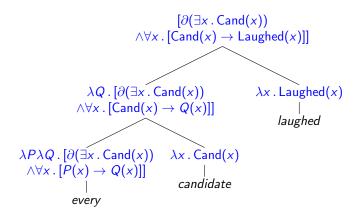
Compositional derivation



Example

```
Every candidate laughed \rightsquigarrow [\partial(\exists x \, . \, \mathsf{Cand}(x)) \land \forall x \, . \, [\mathsf{Cand}(x) \to \mathsf{Laughed}(x)]]
```

Compositional derivation



Comparison with Heim and Kratzer style

$$\lambda Q \cdot [\partial(\exists x \cdot \mathsf{Cand}(x)) \wedge \forall x \cdot [\mathsf{Cand}(x) \to Q(x)]]$$

in Heim and Kratzer style would be:

$$\lambda Q$$
: $\exists x[x \text{ is a candidate }] . \forall x[x \text{ is a candidate } \rightarrow Q(x)]$

The colon and the dot separate the **domain restriction** for the function.

Comparison with Heim and Kratzer style

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Comparison with Heim and Kratzer style

$$[\partial(\exists x \, . \, \mathsf{Cand}(x)) \land \forall x \, . \, [\mathsf{Cand}(x) \to \mathsf{Laughed}(x)]]$$

In Heim and Kratzer style would be:

???

- Barker, Chris. 2005. Remark on Jacobson 1999: Crossover as a local constraint. *Linguistics and Philosophy* 28(4). 447–472.
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